

SETS

- Well defined collection of elements or numbers

- # Types of sets

(1) Null Set :- The set which contains no elements at all is called empty null or void set
Denoted by ' \emptyset ' or '{ }'
ex :- $\{x : x \neq x\}$

(2) Singleton set :- The set which contains only one element is called singleton set ex :- $\{x : x + 5 = 8\}$ or $x = \{3\}$

(3) finite set :- The set which contains finite or limited no. of element is finite set. In other words we can count the different element of a set
ex :- $\{x : 1 \leq x < 7 \text{ and } x \in \mathbb{N}\}$
 $\{2, 3, 4, 5, 6\}$

(4) Infinite set :- The set which is not finite is infinite set ex :- $\mathbb{N}, \mathbb{Z}, \mathbb{Q}$

(5) subset :- Let A and B be two sets

if every element of A is also an element of B then A is called a subset of B and it is denoted by $A \subset B$ and readed as A is subset of B

$$A \subset B \Rightarrow \{x \in B \Rightarrow x \in A\}$$

If A is not a subset of B :- $A \not\subset B$

$$\text{ex:- } A \{2, 4, 6\}; B \{2, 4, 6, 8\}$$

$$\Rightarrow A \not\subset B$$

(6) Proper subset :- Let A and B be two sets then A is called a proper subset if

- A is a subset of B i.e $A \subset B$ and
- there exist at least one element of B which doesn't belongs to A and if is denoted by ' $A \subset B$ ' when A is a proper subset of B then B is called Superset of A and denoted by ' $B \supset A$ '

(7) Improper subset :- Set A is called improper subset of B if and only if $A = B$
Every set is a improper subset of itself or empty set is proper subset of every set $A \neq \emptyset$

(8) Equal set :- Every element of A is also an element of B $\Rightarrow A \subseteq B$
Every element of B is also an element

$A \cup B \subseteq A$

(9) Power set :- Let A be a set of all

subset of A is called the power set

of A it is denoted by $P(A)$

$$\text{ex:- } ① A = \{2, 3\}$$

$$P(A) = \{\emptyset, \{2\}, \{3\}, \{2, 3\}\}$$

$$② A = \{2, 3, 4, 5\}$$

$$P(A) = \{\emptyset, \{2\}, \{3\}, \{4\}, \{5\}, \{2, 3\}, \{2, 4\}, \{2, 5\}, \{3, 4\}, \{3, 5\}, \{4, 5\}, \{2, 3, 4\}, \{2, 3, 5\}, \{3, 4, 5\}, \{2, 3, 4, 5\}\}$$

(10) Index Set :- If $t \in T$ there exist

Set

for ex:- $T = \{a, b, c\}$ then there

are three corresponding set

C_a, C_b, C_c , similarly if $T = \{1, 2, \dots, n\}$ then they are n corresponding set C_1, C_2, \dots, C_n .

(11) Universal Set :-

$$\{x : x \in A_t, t \in T\}$$

$$\text{ex:- If } T = \{1, 2, 3, 4\} \text{ & } A_1 = \{a, b, c, d, e, f\}$$

$$A_2 = \{g, h, i, j, k, l\}, A_3 = \{m, n, o, p, q, r\}$$

$$A_4 = \{s, t, u, v, w, x\}, A_5 = \{y, z, a, b, c, d, e, f, g\}$$

Operations On Set :-

i. Union :- Let A and B be two sets then the union of A & B is the set of all those elements which are in A and B

or in $A \cup B$. Both it is denoted by

' $A \cup B$ ', and read as 'A union B'

$$\text{ex:- } ① A = \{2, 4, 6, 8\}, B = \{12, 13, 14, 15\}$$

$$A \cup B = \{2, 4, 6, 8, 12, 13, 14, 15\}$$

$$② \text{ If } A = \{x : 1 < x < 4, x \in \mathbb{Z}\}, B = \{6 < x < 10, x \in \mathbb{Z}\}$$

then $A \cup B = ?$

$$\Rightarrow A = \{2, 3\}, B = \{7, 8, 9\}$$

$$\Rightarrow A \cup B = \{2, 3, 7, 8, 9\}$$

\Rightarrow Union of Arbitrary collection of sets

Let T be a index set and $t \in T$

then corresponding set A_t then the union of arbitrary collection of sets

at is denoted by $\bigcup_{t \in T} A_t$ and it is

the set of all those elements

which are in A_t for same $t \in T$

$$\{x : x \in A_t, t \in T\}$$

Disjoint sets:- Two set A and B are called mutually disjoint if they have no element common or in other words their intersection is empty set.

Ex:- $A = \{a, b, c\}$, $B = \{e, f, g\}$
then $A \cap B = \emptyset$ hence A & B are disjoint sets.

Difference of two sets

Let A and B be two sets of elements of A and B is the set of elements which belongs to A but doesn't belong to B.

The diff. from B to A is denoted by $A - B$ or $A \sim B$, and read as A difference B or A minus B.

Ex:- $A - B = \{x : x \in A \text{ and } x \notin B\}$

$B - A = \{x : x \in B \text{ and } x \notin A\}$

From the definition we clearly have.

$$\textcircled{Q} \quad A - B \subseteq A$$

$$\textcircled{Q} \quad (A - B) \cap (B - A) = \emptyset$$

$$\begin{aligned} \textcircled{1} \quad (A - B) \cup (B - A) \\ \textcircled{2} \quad (A \cup B) - (A \cap B) \end{aligned}$$

NOTE:- $A \Delta A = \emptyset$
 $A \Delta \emptyset = A$

Complement of set :- Let A be any set and U be a universal set the set $U - A$ is called the complement set of A and it is denoted by A' or A^c or $C(A)$.

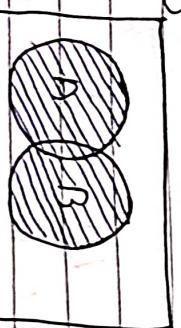
In symbolic notation $A' \text{ or } A^c = \{x : x \in U \text{ and } x \notin A\}$ $U = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$

$$A = \{1, 3, 5, 7, 10\} \text{ then } A' = ?$$

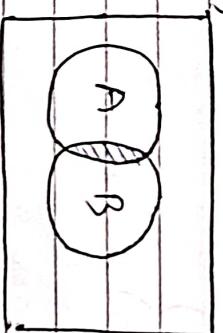
$$A' = \{2, 4, 6, 8, 9\}$$

Venn diagrams

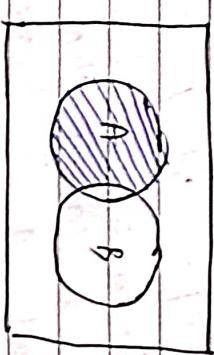
(1) $A \cup B$



(2) $A \cap B$



(3) $A - B$



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Student are either in DM or Java class

$$\text{So } DM = 28, \quad \text{Java} = 30, \quad \text{Both} = 8$$

$$|A \cup B| = |A| + |B| - |A \cap B|$$

$$|A \cup B| = 28 + 30 - 8$$

$$|A \cup B| = 50$$

$$(2) |A_1 \cup A_2 \cup A_3| = |A_1| + |A_2| + |A_3| - |A_1 \cap A_2| - |A_1 \cap A_3| - |A_2 \cap A_3| + |A_1 \cap A_2 \cap A_3|$$

Q2) A shop of computer has 8 computers having the following specification

Let A_1 A_2 A_3

Computer floating pt Magnetic Graphic
with arithmetic unit disk display

I Y Y

II Y Y

III N Y

IV N Y

V Y N

VI N Y

VII N N

VIII Y N

Principle of Inclusion & Exclusion

If A and B are infinite set then

$$(1) |A \cup B| = |A| + |B| - |A \cap B|$$

Q1) At IIPS there are 28 students in Discrete mathematics class

30 students in Java & students in both class How many

find the no. of computer hardware one or more kind of hardware?

$$\text{Soln} \quad A_1 = 3, A_2 = 5, A_3 = 4, \\ A_1 \cap A_2 = 2, A_1 \cap A_3 = 2, A_2 \cap A_3 = 1, \\ A_1 \cap A_2 \cap A_3 = 1$$

$$|A_1 \cup A_2 \cup A_3| = |A_1| + |A_2| + |A_3| - |A_1 \cap A_2| - |A_1 \cap A_3| - |A_2 \cap A_3| + |A_1 \cap A_2 \cap A_3| \\ = 3 + 5 + 4 - 2 - 2 - 1 + 1 \\ = 11$$

(Q4) In High school there are 55 students in either algebra, biology or chem class. 28 student in algebra class, 30 student in bio, 29 student in both algebra and bio & student in both algebra and chem, 16 student in both bio & chem, and 5 student in both algebra & chem.

$$\text{Soln} \quad 55 = 28 + 30 + 24 - 8 - 16 - 5 \\ = 82 - 29 + (A \cap B \cap C) \\ (A \cap B \cap C) = 21$$

Ordered pair of Cartesian Product
An ordered pair consist of two objects in a given fixed order note that an ordered pair is not set consisting of two elements the ordering of the two objects is imp: the two objects need not to be distinct we can denote ordered pair by (x, y)

$$(B \cup U) = 21 + 26 + 29 - 14 - 15 - 12 + 8 \\ = 82 - 41 \\ = 41$$

\Rightarrow Cartesian Product
Let A and B be any two sets the set of all ordered pairs such that first member of ordered pair is an

element of $A \times B$ & the second part member of B is called the Cartesian product of $A \times B$ and is written as $A \times B$

$A \times B \times C$ consists of all ordered triplets (a, b, c) where $a \in A$, $b \in B$, etc. These elements of $A \times B \times C$ can be systematically obtained by a so it is called tree diagram.

$$A \times B = \{(x, y) : x \in A \text{ and } y \in B\}$$

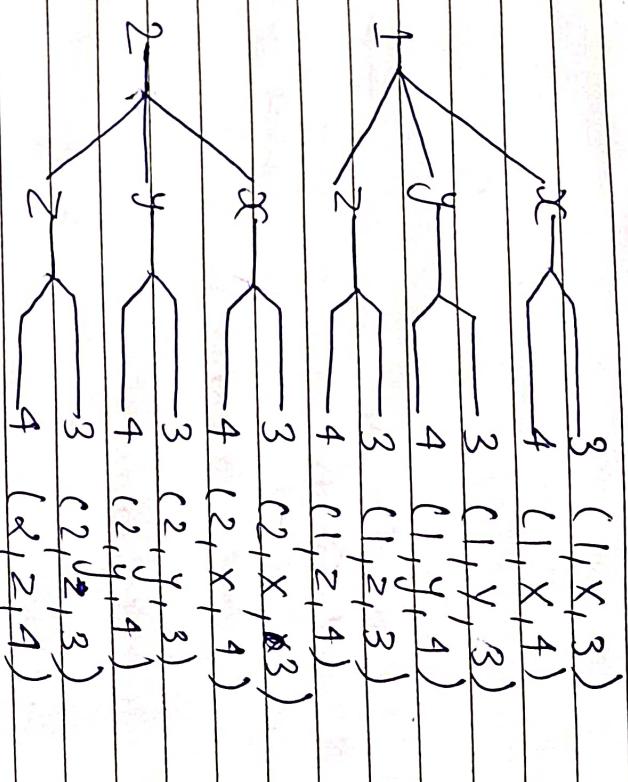
$$B \times A = \{(x, y) : x \in B \text{ and } y \in A\}$$

(Q) If $A = \{1, 2, 3\}$ and $B = \{1, 2, 3\}$
then what is $A \times B$, $B \times A$, $A \times A$, $B \times B$

Soln
 $\Rightarrow A \times B = \{(1, 1), (1, 2), (1, 3), (2, 1), (2, 2), (2, 3), (3, 1), (3, 2), (3, 3)\}$

$$\Rightarrow B \times A = \{(1, 1), (1, 2), (1, 3), (2, 1), (2, 2), (2, 3), (3, 1), (3, 2), (3, 3)\}$$

$$\Rightarrow A \times A = \{(1, 1), (1, 2), (1, 3), (2, 1), (2, 2), (2, 3), (3, 1), (3, 2), (3, 3)\}$$



The element of $A \times B \times C$ are 12 ordered triplets to right side of tree diagram.

$$n(A) = 2$$

$$n(B) = 3$$

$$n(C) = 3$$

Q) $A = \{1, 2\}$, $B = \{1, 2, 3\}$, $C = \{3, 4\}$

find $A \times B \times C$ and $n(A \times B \times C)$.

Soln. $A \times B \times C = \{(1, 1, 3), (1, 1, 4), (1, 2, 3), (1, 2, 4), (2, 1, 3), (2, 1, 4), (2, 2, 3), (2, 2, 4)\}$

$$A \times 1$$

Q) Let $A = \{a, b, c\}$, $B = \{b, c, d\}$
and $C = \{a, d\}$. Construct a
true diagram of $A \times B \times C$

$$\begin{array}{c} B \\ \downarrow \\ R \\ \downarrow \\ A \end{array}$$

Now suppose a relation from set A to set B is expressed by the statement is less than or equal to.

$$R = \{(1, 2), (1, 4), (1, 3), (3, 4), (3, 8), (5, 8)\}$$

Let A and B be two sets a relation from A to B is the subset of $A \times B$. ($R \subseteq A \times B$) and it is denoted by ' \rightarrow ' R . R is a relation from A to B $\Rightarrow R \subseteq A \times B$.

$$R = \{(x, y) : x \leq y \text{ where } x \in A \text{ and } y \in B\}$$

\Rightarrow Total no. of different relations of the set A has m elements and the set B has n elements then the total no. of relation from set A to set B is 2^{mn} since $A \times B$ has $m \times n$ elements in all.

Relation

In order to express a relation from set A to set B we always \Rightarrow domain of range of a relation induced a statement which connects the element of A with the elements of B i.e. $R \subseteq A \times B$. The domain

R is the set of all first element of the ordered pair which belongs to R and the range of R is the set of all second element of the ordered pair.

→ Domain of R = set of first elements of the ordered pair which belongs to R . $= \{x : x \in A\}$ and $(x, y) \in R$ for some $y \in B\}$

→ Range of R = set of all second elements of the ordered pair which belongs to R = $\{y : y \in B\}$ and $(x, y) \in R$ for some $x \in A\}$

⇒ Inverse Relation

If R is a relation from set A to B then the inverse of R is a relation from B to A and it is denoted by $R^{-1} = \{(y, x) : (x, y) \in R\}$

Thus to find R^{-1} we write in reverse order all ordered pairs belonging to R .

Range of $R^{-1} \Rightarrow$ Domain of R

- ① Reflexive :- If R is a relation in the set A then R is called reflexive

⇒ The composition relation of two relations

Let R , S be three sets. Suppose R is a relation from set A to B and S is the relation from set B to C . The composition relation of two relations R & S is the relation from set A to C and it is denoted by SOR and defined as $SOR = \{(a, c) : \text{there exists an element } b \in B \text{ such that } (a, b) \in R \text{ and } (b, c) \in S\}$ where $a \in A$, $c \in C$. Hence $(a, b) \in R$, $(b, c) \in S \Rightarrow (a, c) \in SOR$

⇒ Binary Relation OR Relation on a Set

Suppose A is a non empty set of relation R in set A is the subset of $A \times A$ clearly both the co-ordinates of ordered pairs in R are the elements of the set A the relation R in the set A is called a binary relationship

Properties of Binary Relation

relation if every element of A is related to itself

$(q, q) \in R$, $\forall q \in A$

$qRa \Leftrightarrow q, \in A$

$R \subseteq A \times A$

if $A = \{1, 2, 3\}$
 $R = \{(1, 1), (2, 2), (3, 3)\}$

\Rightarrow Symmetric Relation

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If R is a relation in the set A then R is called symmetric relation if a is R related to b

i.e. $(a, b) \in R \Rightarrow (b, a) \in R$

i.e. $aRb \Rightarrow bRa$

where $a, b \in A$

ex:- If $A = \{2, 4, 5, 6\}$

Symmetric $R_1 = \{(2, 4), (4, 2), (4, 5), (5, 4)\}$

Not $\textcircled{B} R_2 = \{(2, 4), (2, 6), (6, 2), (5, 4), (4, 5)\}$

ex:- If $A = \{2, 4, 6\}$

$R = \{(2, 2), (4, 4), (6, 6)\}$

\Rightarrow Antisymmetric Relation & If R is a relation in the set A then R is called antisymmetric

Antisymmetric if $(q, b) \in R$ and $(b, q) \in R$

$\Rightarrow Q = b$ where $q, b \in A$

ex:- in a set of natural no. the relation "a divides b" is antisymmetric

since "a divides b" and "b divides a" is possible only when $a = b$ it means if

the given relation be denoted by R then $(q, b) \in R$ and $(b, q) \in R \Rightarrow q = b$

\Rightarrow Transitive relation :- If R is a relation in the set A then R is called transitive

relation if a is R related to B and b is R related to C then there must

be a is R related to C

ex:- $A = \{1, 2, 3\}$

then $R = \{(1, 2), (2, 3), (1, 3)\}$

$\Rightarrow (q, b) \in R$ and $(b, c) \in R \Rightarrow (q, c) \in R$

$q, b, c \in A$

\Rightarrow Identity relation :- If R is a relation in set A then R is called identity

relation if a is R related to b by $aRb \Rightarrow a = b$ $\forall q, b \in A$

ex:- If $A = \{2, 4, 6\}$

$R = \{(2, 2), (4, 4), (6, 6)\}$

Note:- every identity relationship is reflexive but the converse is not true.

\Rightarrow Equivalence relation:- Suppose R is the relation in set A then R is

called equivalence relation if

R is reflexive $(q, q) \in R$ $\forall q \in A$

R is symmetric $(q, b) \in R \Rightarrow (b, q) \in R$

where $a, b \in A$

(3) R is transitive $(a, b) \in R$ and $(b, c) \in R$

then $(a, c) \in R$ also where $a, b, c \in A$

(Ans) If $R_1 = \{(a, b) \mid a - b \text{ is an integer}\}$

$R_2 = \{(a, b) \mid a - b \text{ is divisible by } 3\}$

$R_3 = \{(a, b) \mid a - b \text{ is an odd no.}\}$

$R_4 = \{(a, b) \mid a - b \text{ is an even no.}\}$

Sol: $R_1 = \{(a, b) \mid a - b \text{ is an integer}\}$

i) Reflexive :- Let $2 \in \mathbb{Z}$ belongs to the set

then, $2 - 2 = 0$ it is an even integer

Hence it is reflexive relation

ii) Symmetric :- Let $1, 2 \in \mathbb{Z}$

then, $1 - 2 = -1 \quad 2 - 1 = 1$

\Rightarrow Hence it is symmetric relation

Both are integers hence it is symmetric

iii) Transitive :- Let $1, 2, 3 \in \mathbb{Z}$

then $(1, 2) \in R$ and $(2, 3) \in R$ then $(1, 3)$

also $\in R$ $\Rightarrow 1 - 1 = 0$, 0 is not an odd no.

$\Rightarrow 1 - 2 = -1 \quad 2 - 3 = -1 \quad 1 - 3 = -2$

\Rightarrow all three are integers

Hence it is transitive and it is equivalence relation

Q) $R_2 = \{(a, b) \mid a - b \text{ is divisible by } 3\}$

i) Reflexive :- let $1 \in \mathbb{Z}$

$\Rightarrow 1 - 1 = 0 \Rightarrow 0$ is divisible by 3

\Rightarrow Hence it is reflexive

ii) Symmetric :- let $3 \in \mathbb{Z}$

$\Rightarrow 3 - 3 = 0$

\Rightarrow Hence it is symmetric

iii) Transitive :- let $3, 6 \in \mathbb{Z}$

$\Rightarrow (3, 6) \in R$ and $(6, 9) \in R$

$\Rightarrow (3, 9) \in R$

$\Rightarrow 3 - 6 = -3 \quad 6 - 9 = -3 \quad 3 - 9 = -6$

$\Rightarrow -3$ is divisible by 3

$\Rightarrow -3$ is divisible by 3

$\Rightarrow -6$ is divisible by 3

\Rightarrow Hence it is transitive relation also

Q) $R_3 = \{(a, b) \mid a - b \text{ is an odd no.}\}$

i) Reflexive :- let $1 \in \mathbb{Z}$

$\Rightarrow 1 - 1 = 0$

\Rightarrow Hence it is not reflexive

ii) Symmetric :- let $-2 \in \mathbb{Z}$

$\Rightarrow -2 + 2 = 0$

\Rightarrow Hence it is symmetric

iii) Transitive :- let $1, 2, 3 \in \mathbb{Z}$

$\Rightarrow (1, 2) \in R$ and $(2, 3) \in R$

$\Rightarrow (1, 3) \in R$

$\Rightarrow 1 - 2 = -1 \quad 2 - 3 = -1 \quad 1 - 3 = -2$

\Rightarrow Hence it is transitive and it is symmetric

Q) $R_4 = \{(a, b) \mid a - b \text{ is an even no.}\}$

i) Reflexive :- let $1 \in \mathbb{Z}$

$\Rightarrow 1 - 1 = 0$

\Rightarrow Hence it is reflexive

ii) Symmetric :- let $2 \in \mathbb{Z}$

$\Rightarrow 2 - 2 = 0$

\Rightarrow Hence it is symmetric

iii) Transitive :- let $1, 2, 3 \in \mathbb{Z}$

$\Rightarrow (1, 2) \in R$ and $(2, 3) \in R$

$\Rightarrow (1, 3) \in R$

$\Rightarrow 1 - 2 = -1 \quad 2 - 3 = -1 \quad 1 - 3 = -2$

\Rightarrow Hence it is transitive and it is symmetric

Hence it is not transitive
 \Rightarrow It is not equivalence relation also

compute any one of them

$$\text{Q) } R_4 = \{(9, 6) | 9-6 \text{ is an even no.}\}$$

Soln i) reflexive:- let $1 \in \text{set}$

$$\Rightarrow 1-1=0, 0 \text{ is an even no.}$$

Hence it is reflexive

ii) symmetric:- let $1, 3 \in \text{set}$

$$\Rightarrow (1, 3) \in R_4 \Rightarrow (3, 1) \in \text{set}$$

$$1-3=-2 \quad 3-1=2$$

\Rightarrow R_4 is an even no. $\Rightarrow R_4$ is an even no.

Hence it is symmetric

iii) transitive:- let $1, 3, \text{ and } 5 \in \text{set}$

$$\Rightarrow (1, 3) \in R_4 \text{ and } (3, 5) \in R_4 \text{ then}$$

$$(1, 5) \text{ also } \in R_4$$

$$\Rightarrow 1-3=-2 \quad 3-5=-2 \quad 1-5=-4$$

$$\Rightarrow -2 \text{ is an even no.} \Rightarrow -2 \text{ is even} \Rightarrow -4 \text{ is even}$$

Hence it is transitive and it is

is also equivalence.

Equivalence Class

$$[x] = \{y | y \in A \text{ and } (x, y) \in R\}$$

$$A = \{1, 2, 3, 4, 5\}$$

$$R = \{(1, 1), (2, 2), (3, 3), (4, 4), (5, 5), (1, 2), (2, 1), (3, 1), (1, 3), (3, 2)\}$$

$$\begin{aligned} [1] &= \{1, 2\} \\ [2] &= \{2, 1\} \text{ and } \{1, 2\} \text{ are same} \\ [3] &= \{3\} \\ [4] &= \{4, 5\} \\ [5] &= \{5, 4\} \end{aligned}$$

$$\begin{aligned} P_1 &= \{1, 2\}, P_2 = \{3\}, P_3 = \{4, 5\} \\ P_1 \cup P_2 \cup P_3 &= A \\ \Rightarrow P_1 \cap P_2 \cap P_3 &= \{\phi\} \end{aligned}$$

25/10/2023

Representing Relation using matrices
 A relation b/w finite sets can be represented using a zero-one matrix.

Suppose that R is the relation from $A = \{a_1, \dots, a_n\}$ to $B = \{b_1, \dots, b_m\}$. The relation R can be represented by the matrix $M_R = [m_{ij}]$ where $m_{ij} = 1$, if $(a_i, b_j) \in R$

Suppose that $A = \{1, 2, 3\}$ and $B = \{1, 2\}$
 Let R be the relation from $\text{set } A$ to set B containing (a, b) if $a \in A, b \in B$ and $a > b$. Then what is the matrix representation of R ?

Hence it is transitive.

for ex:-

(i) Reflexive :- for each $x \in I$ we have
 $x - x = 0$ is divisible by 5
 $\Rightarrow xRx$

Thus $\forall x \in I$ we get xRx
 Therefore R is reflexive

relation $R = \{(x, y) : x, y \in N \text{ and}$

(i) let I be the set of all integers
 Let's define the relation R in
 set I such that xRy if $x-y$ is
 divided by 5. $x \in I$ and $y \in I$
 i.e $R = \{(x, y) : x \in I, y \in I \text{ and}$
 $x-y \text{ divisible by } 5\}$

Soln (i) Reflexive :- let $I \subseteq I$
 $\Rightarrow 1 - 1 = 0$

$\Rightarrow 0$ is divisible by 5
 Hence it is reflexive.

(ii) Symmetric :- let $x, y \in I$ we have
 $x-y$ is divisible by 5
 $\Rightarrow -(x-y)$ is divisible by 5
 $\Rightarrow (y-x)$ is divisible by 5
 $\Rightarrow y-x$ is divisible by 5
 Thus $xRy \Rightarrow yRx$

Therefore R is symmetric

(iii) Transitive :- let $x, y, z \in I$
 we have xRy, yRz

$\Rightarrow x-y$ divisible by 5
 $\Rightarrow y-z$ divisible by 5
 $\Rightarrow x-z$ divisible by 5

(iv) Symmetric :- let $I_0, I_1, I_2 \subseteq I$
 $\Rightarrow s - 10 = -5 \mid 10 - s = s$. $\text{if } R_2 = xR_2 y \text{ then } f(I_2)$
 $\Rightarrow -5$ is divisible by 5.

Hence it is symmetric.

Since R is reflexive, symmetric &
 Transitive $\Rightarrow R$ is equivalence relation

(v) Transitive :- let $s_1, 10, 15 \in I$.

$\Rightarrow (s_1, 10) \in R$ and $(10, 15) \in R$

$\Rightarrow (s_1, 15) \in R$

$\Rightarrow s_1 - 10 = -s \mid 10 - 15 = -s \mid s_1 - 15 = -s$

relation are equivalence relation

$$R_1 = \{(x, y) : x, y \in L\}$$

Soln (i) reflexive for any $x \in L$ we have $x \parallel x$ to itself

$$\Rightarrow (\underline{x}, \underline{x}) \in R \Rightarrow x \cdot R \cdot x \Rightarrow R_1$$

Hence R_1 is reflexive

(ii) symmetric:- for any $x \in L$ and $y \in L$ we have $x \parallel y$

$$\Rightarrow y \parallel x$$

$$\Rightarrow x \cdot R \cdot y \Rightarrow R_1$$

Hence R_1 is symmetric

Soln Let R be an equivalence relation on $N \times N$ where N is a set of the integers defined by $(a, b) R (c, d) \Rightarrow a+d = b+c$ and $a, b, c, d \in N$ then prove that R is an equivalence relation or not

Partial ordered relation

A relation R on a set A is called Partial ordered relation if R is

① reflexive i.e. $aRa \forall a \in A$

② R is anti symmetric i.e. $aRb, bRa \Rightarrow a=b$ for each $a, b \in A$

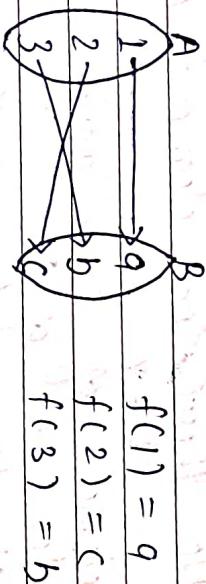
③ R is transitive i.e. $aRb, bRc \Rightarrow aRc$ for each $a, b, c \in A$.

function:- A function is a special kind of relation. In many instances we assigned each element of first set to a particular element of second set for example suppose that each student in discrete mathematics class is assigned a letter from the set $A = \{a, b, c, d, e\}$ and suppose that these are

$a = \text{Outstanding}$, $b = \text{Excellent}$,
 $c = \text{Very Good}$, $d = \text{Good}$, $e = \text{Fair}$

Suppose that to each element of A there is assigned a unique element of B . the collection of such assignment is called a function or Mapping or Map from A to B if there are two sets. We denote a function f from A to B by $f: A \rightarrow B$ and read as f of A onto B .

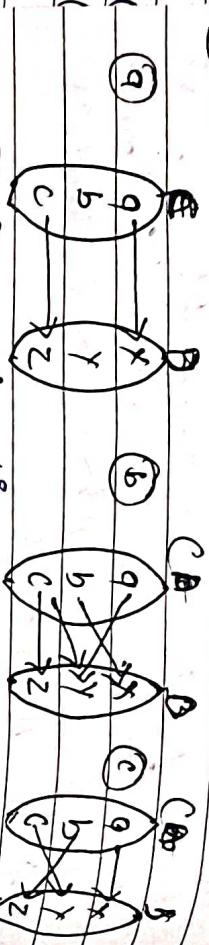
Where the set A is a domain of f and the set B is the co-domain of f we write $f(a)$ and read as f of a for the element of B that f assigns to $a \in A$ and it is called the value of f at ' a ' or the image of ' a ' under f .



The concept of function is extremely important in discrete mathematics. The functions are used in defining such as discrete structures as sequence of strings to be equal if $f: A \rightarrow B$ and $g: A \rightarrow B$ are defined to be equal if $f(g)(a) = g(a) \forall a \in A$. The negation of $f = g$ is $f \neq g$ and in the statement of $f \neq g$ there exists an $a \in A$ for which $f(a) \neq g(a)$.

(Ans) state whether or not each diagram below defined as function from

$$A = \{a, b, c\} \quad D = \{x_1, y_1, z_1\}$$



Soln C is a function



Soln No there no element in D

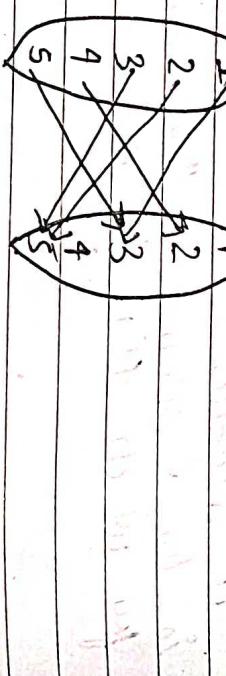
$$b \in A$$

element of C is assigned a unique element of D

No two element x and y $\in C$

$$(3) f(3) = ?$$

Ques) Consider the set $A = \{1, 2, 3, 4, 5\}$ the function $f : A \times A$ defined by



so (1) $f(1) = 3$, $f(2) = 5$, $f(3) = 5$, $f(4) = 2$

The arrow indicates the image of an element.

(2) The image $f(A)$ of f consist all the image value. Only 2, 3, 5 appears as a image of any element
Thus $f(A) = \{2, 3, 5\}$

(i) find $f(T)$ where $T = \{1, 2\}$

$$(2) f^{-1}(T) = \{3, 5\}$$

$$(1) f^{-1}(3) = \{1, 2\}$$

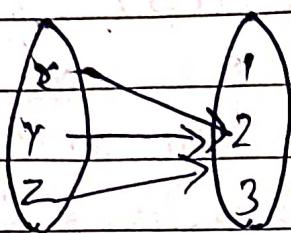
$$\{3, 5\} = \{3, 5\}$$

$$(2) f^{-1}(2) = \{1, 2\}$$

$$(3) f^{-1}(5) = \{1, 2\}$$

$$(4) f^{-1}(3) = \{1, 2\}$$

element this can be represented by diagram where all images



$$f(x) = 2$$

$$f(y) = 2$$

$$f(z) = 3$$

Types of Mapping

→ Injective or one-one mapping

A mapping ~~$f(x)$~~ f of X into Y is said to be injective or one-one mapping if distinct element of X have distinct images in Y . It is called injective or one-one mapping.

e.g. $f: X \rightarrow Y$ is a injective (one-one) mapping if and only if

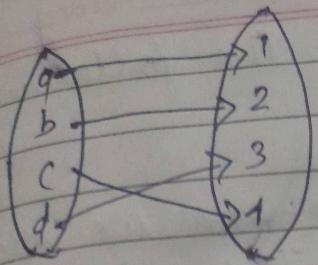
$$f(x_1) = f(x_2) \Rightarrow x_1 = x_2$$

ex:- if $X = \{a, b, c, d\}$, $Y = \{1, 2, 3, 4\}$

and $f: X \rightarrow Y$ defined as

$$f(a) = 1, f(b) = 2, f(c) = 4$$

$$f(d) = 3$$



hence it is
injective or
one-one

→ Surjective or Onto mapping :-

If the mapping f of X into Y is such that every element of Y has the image of atleast one element of X then the mapping is called surjective or onto.

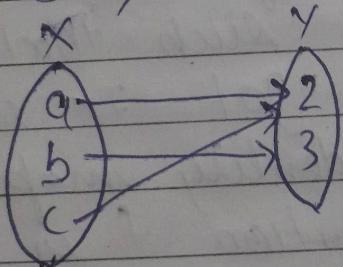
In other words

$f: X \rightarrow Y$ is onto \bar{Y} if given $y \in Y$ there exist an element $x \in X$ such that $y = f(x)$

ex: if $X = \{a, b, c\}$, $Y = \{2, 3\}$ and if

$f: X \rightarrow Y$ is defined by $f(a) = 2$

$f(b) = 3$, $f(c) = 2$



It is surjective.

→ Bijective or One-One Onto mapping

A mapping which is one-on-one as well as onto is called bijective or one-one onto mapping. To determine whether the mapping is bijective.

we have to follow the following procedure :-

- To show if f is one-one we must show that $f(x_1) = f(x_2) \Rightarrow x_1 = x_2$
- To show if f is onto y we must show that for each $y \in Y$ there exist an element $x \in X$ such that $f(x) = y$

ex:- If $X = \{a, b, c, d\}$ $Y = \{1, 2, 3, 4\}$
and $f(a) = 1$, $f(b) = 2$, $f(c) = 4$
 $f(d) = 3$

It is one-one and onto i.e. surjective

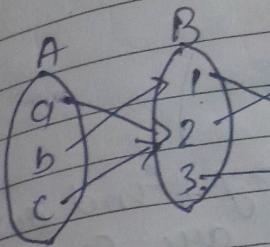
Invertible function :-

function f of X into Y is said to be invertible if there exist a function $g: Y \rightarrow X$ such that

$f \circ g = I_Y$ and $g \circ f = I_X$ where I_X and I_Y are identity map in such case, a function f is called inverse function and it is denoted by ' f^{-1} '.

Let the function $f: A \rightarrow B$, $g: B \rightarrow C$ and $h: C \rightarrow D$ be defined by below diagram. Determine which of the function are

① one-one ② on



② $\Rightarrow g: f: A$
because every image
 $\Rightarrow g: B \rightarrow C$
 $\Rightarrow h: C \rightarrow D$
 $h(x)$

② $f: A \rightarrow B$

3 does
 $\Rightarrow g: B \rightarrow C$
2 does
 $\Rightarrow h: C \rightarrow D$

③ $f: A \rightarrow B$

① one-one ② onto ③ Invertible

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following

we must
 $\Rightarrow x_1 = x_2$

we must
 $\in Y$ there
 such that

$\{1, 2, 3, 4\}$
 $f(c) = 4$

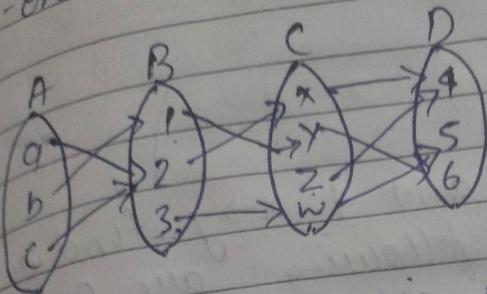
to i.e. ~~Be~~ surjective

is said
 here exist g
 that

$\exists x$ where $\exists x$

map in
 f is called
 t is denoted

$\rightarrow B$
 $g: B \rightarrow X$
 need by below
 which of the



① $\Rightarrow g: f: A \rightarrow B$ is not one-one
 because $f(a) = 2$ & $f(c) = 2$

$\Rightarrow g: B \rightarrow C$ is one-one because
 every element of B has a unique
 image

$\Rightarrow h: C \rightarrow D$ is not one-one because
 $h(x) = 4$ and $h(z) = 4$

② $f: A \rightarrow B$ is not onto because
 3 doesn't have any image in A

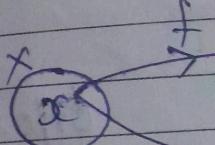
$\Rightarrow g: B \rightarrow C$ is not onto because
 z doesn't have any image in B

$\Rightarrow h: C \rightarrow D$ is onto as every element
 of D has atleast one image in C

③ $f: A \rightarrow B$

Composite of f
Let there be
Suppose two +
defined by $y \in Y$ and $g =$
 $z = g(y)$

f and g is
by $f: X \rightarrow Y$
 $gof: X \rightarrow Z$
 $(gof)(x)$



Ques) If $X = \{0, \infty\}$ and $f: X \rightarrow Y$ find out which of the following are onto (surjective) mapping

i) $f(x) = x^2$ ii) $f(x) = x^3 + 1$
iii) $f(x) = x + 2$

Soln We have $f(x)$ when $0 \in X$

i) $\Rightarrow f(0) = 0^2 = 0 \in X$

when $\infty \in X$

$f(\infty) = \infty^2 \Rightarrow \infty \in X$

Hence it is onto

ii) $f(x) = x^3 + 1$

$\Rightarrow f(0) = 0 + 1 \Rightarrow 1 \notin X$

$\Rightarrow f(\infty) = \infty + 1 \Rightarrow \infty \in X$

Since $1 \notin X$ so it is not onto

iii) $f(x) = x + 2$

$\Rightarrow f(0) = 0 + 2 \Rightarrow 2 \notin X$

$\Rightarrow f(\infty) = \infty + 2 \Rightarrow \infty \in X$

Since $2 \notin X$ so it is not onto

Ex:- Let f a
set of
integers
 $g(x) = 3x$
what is
what is
som $(fog)x$

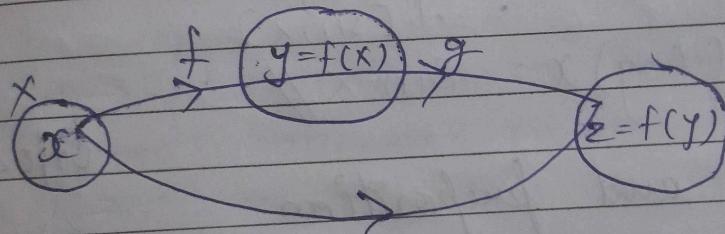
$2(3x+2)$
 $(gof)x$

Composite of function

let there be 3 sets $X, Y \& Z$

suppose two functions $f: X \rightarrow Y$ is defined by $y = f(x)$ where $x \in X$ & $y \in Y$ and $g: Y \rightarrow Z$ is defined by $z = g(y)$ where $y \in Y$ and $z \in Z$

f and g is composite mapping denoted by $g \circ f$ and if a function $f: X \rightarrow Y$ is denoted by $g \circ f: X \rightarrow Z$ where $(g \circ f)(x) = g[f(x)] \forall x \in X$ *



ex:- let f and g be the function from set of integers to the set of integers defined by $f(x) = 2x+3$ and $g(x) = 3x+2$

what is the composition of f and g and what is the composition of g and f

$$(f \circ g)x = f[g(x)] = f(3x+2)$$

$$2(3x+2)+3 = 6x+4+3 = 6x+7 \quad (1)$$

$$(g \circ f)x = g[f(x)] = g[2x+3]$$

$$3(2x+3) + 2(8x+11) \quad \text{--- (2)}$$

Ques1) If $f: \mathbb{R} \rightarrow \mathbb{R}$ defined by $f(x) = x^2$ & $x \in \mathbb{R}$ and $g: \mathbb{R} \rightarrow \mathbb{R}$ is defined by $g(x) = \sin x$ of $x \in \mathbb{R}$ then find gof and fog and also show whether $gof = fog$ or $gof \neq fog$

Soln

$$\Rightarrow (gof)(x) = g[f(x)] \\ (gof)(x) = g[\sin x] //$$

$$\Rightarrow (fog)x = f[g(x)] \\ = f[\sin x]$$

$$\Rightarrow (fog)x = \sin^2 x //$$

07/02/2023

Logic and preposition

\Rightarrow Sentence :- A group of words which may not have its value true or false

\Rightarrow Statement :- It is a special kind of sentence which is defined as its value true or false.

\Rightarrow Logical connections :- They are the words or symbols used to combine two

- Sentences to
or a comp
- 1 connective words
 - 1 not
 - 2 and
 - 3 or
 - 4 if...then
 - 5 iff or 'if & only if'

① NEGATION

Truth

Inpu

O(F

I(T

x -

② Conjunction

statements

word are

word 'A

represent

P and

AND

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sentences to form a compound sentence
 by a compound statement.

1 $f(x) = x^2$
 is defined
 if & then
 and also
 fog by

Connective words	Name of connectives	Symbol	Rank
not	Negation	\sim or \neg	1
and	Conjunction	\wedge	2
or	Disjunction	\vee	3
if...then	Conditional	\Rightarrow or \rightarrow	4
iff or 'if & only if	Bi-conditional	\Leftrightarrow or \leftrightarrow	5

① NEGATION :-

Truth Table		$y = \neg x$
Input	Output	
0 (F)	1 (T)	
1 (T)	0 (F)	

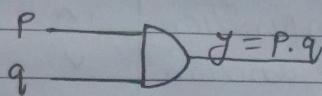
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$$x \longrightarrow \neg y = \neg \neg x$$

② Conjunction :- Suppose there are two statements P and q when two statement word are combined while using the word 'AND' then we get a new statement represented as : $P \wedge q$ and read as P and q .
 $AND \rightarrow P \cdot q$

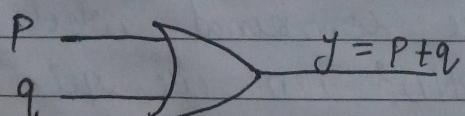
ds which
 true or
 kind of
 as its
 the words
 kine two

Input		Output
P	q	$P \wedge q$
T	T	T
T	F	F
F	T	F
F	F	F



⑤ Disjunction :- When two statement p and q are join by using the word 'or' then we get a new statement represented by $P \vee q$ and read as P or q , its value is true if any one of the statement is true.

Input		Output
P	q	$P \vee q$
T	T	T
T	F	T
F	T	T
F	F	F



OR $\rightarrow P + q$

⑥ Conditional :-
-ent then the
P then q is
ment and is
Its value is
 $\delta = q$ is false

Input	
P	q
T	T
T	F
F	T
F	F

⑥ Bi-condition statements
the type
bi-condi
value is
are either

Input	
P	q
T	T
T	F
F	T
F	F

④ Conditional :- If $P \& q$ are two statement then the statement of type if P then q is called conditional statement and is represented by $P \rightarrow q$. Its value is false if P is true & q is false.

Input		Output
P	q	$P \rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

⑤ Bi-Conditional :- If P and q are two statements then the statement of the type $P \rightarrow q$ and $q \rightarrow P$ is called bi-conditional statement and its value is true when both $P \& q$ are either true or false.

Input		$P \rightarrow q$	$q \rightarrow P$	Output
P	q	$P \rightarrow q$	$q \rightarrow P$	$P \rightarrow q$ and $q \rightarrow P$
T	T	T	T	T
T	F	F	T	F
F	T	T	F	F
F	F	T	T	T

Tautology :- It is a proposition which is true for all truth values of its sub propositions on components.

It is also called as logically valid or logically true statements. Any compound statement is said to be a tautology if it has all the truth values true.

Contradiction :- A contradiction is a proposition which is always false for all truth values of its sub propositions or components.

Also called as logically invalid or logically false statement. Any compound statement is said to be a contradiction if it has all the truth values false.

Logical equivalence :- Two statements are said to be logically equivalent if truth values of both the statements are always identical.

And it is written as logically equivalent q { $P = q$ }

Contingency :- ex:- if P & q then find whether

- ① $P \wedge q \Rightarrow (P \vee q)$
- ② $(P \wedge q) \Leftrightarrow (P \rightarrow q)$
- ③ $(P \leftarrow q) \Leftrightarrow (P \rightarrow q)$
- ④ $(P \rightarrow q) \wedge (q \rightarrow r) \Rightarrow (P \rightarrow r)$

Ans	P	q
	T	T
	T	F
	F	T
	F	F

yes tautology

②	P	q
	T	T
	T	F
	F	T
	F	F

yes tautology

③	P	q
	T	T
	T	F
	F	T
	F	F

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Position
truth
position

as logically
elements
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statements
equivalent
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(contingency:-
ex:- if $P \& q$ are two statements
then find whether :-

$$\textcircled{1} P \wedge q \Rightarrow P \text{ is tautology or not}$$

$$\textcircled{2} (P \wedge q) \Rightarrow (P \vee q)$$

$$\textcircled{3} (P \Leftrightarrow q) \Leftrightarrow (P \Rightarrow q) \wedge (q \Rightarrow P)$$

$$\textcircled{4} (P \Rightarrow q) \wedge (q \Leftarrow r) \rightarrow (P \Rightarrow r)$$

Ans	P	q	$P \wedge q$	$P \wedge q \Rightarrow q$
	T	T	T	T
	T	F	F	T
	F	T	F	T
	F	F	F	T

yes tautology

②	P	q	$P \wedge q$	$P \vee q$	$(P \wedge q) \Rightarrow (P \vee q)$
	T	T	T	T	T
	T	F	F	T	T
	F	T	F	T	T
	F	F	F	F	T

yes tautology

③	P	q	$P \Leftrightarrow q$	$P \Rightarrow q$	$q \Rightarrow P$	$(P \Rightarrow q) \Leftrightarrow (q \Rightarrow P)$
	T	T	T	T	T	T
	T	F	F	F	T	T
	F	T	F	T	F	F
	F	F	T	T	T	T

T	T	T	T	T	T	T
T	F	F	F	T	F	T
F	T	F	T	T	F	F
F	F	T	T	T	T	T

Output

T
T
F
T

No tautology

④

$(P \vee q) \wedge (\sim P) \wedge (\sim q)$
F
F
F
F

yes it is

P	q	r
T	T	T
T	T	F
T	F	T
T	F	F
F	T	T
F	T	F
F	F	T
F	F	F

Ques) PT $(P \vee q) \wedge (\sim P) \wedge (\sim q)$ is contradiction or not
Please that $P \Rightarrow (q \Rightarrow r) \equiv (P \wedge q) \Rightarrow r$
whether the above preposition logical equivalence or not.

P	q	$\sim P$	$\sim q$	$P \vee q$	$(P \vee q) \wedge (\sim P) \wedge (\sim q)$
T	T	F	F	T	F
T	F	F	T	T	F
F	T	T	F	T	T
F	F	T	T	F	F

$$(P \vee q) \wedge (\neg P) \wedge (\neg q)$$

F

F

F

F

yes it is a contradiction

P	q	r	(q \Rightarrow r)	$P \Rightarrow (q \Rightarrow r)$	(P \wedge q)
T	T	T	T	T	T
T	T	F	F	F	T
T	F	T	T	T	F
T	F	F	F	F	F
F	T	T	T	T	F
F	T	F	F	F	F
F	F	T	T	T	F
F	F	F	F	F	F

$$P \wedge q \Rightarrow r$$

m

$$q \wedge (\neg P)$$

F

F

T

F

Boolean Algebra

Let B is the set (non empty) with element $\{a, b, c\}$ which is defined on two operation '+' and '.', then the algebraic structure ($P+P$) over B is said to be Boolean algebra if B satisfied the following properties or laws or proposition

→ closure law :-

$$\begin{aligned} \text{if } a, b \in B &\Rightarrow a+b \in B \\ &\Rightarrow ab \in B \end{aligned}$$

→ commutative law :-

$$\begin{aligned} a+b &= b+a \\ a.b &= b.a \quad \forall a, b \in B \end{aligned}$$

→ Distributive law :-

$$a+(b.c) = (a+b)(a+c) \quad \forall a, b, c \in B$$

$$a.(b+c) = ab+ac \quad \forall a, b, c \in B$$

→ Identity law :- there should be identity elements on both side of operation identity element of $+ = 0$

$$a+I = a \quad I = 0$$

for multiplication
 $- I \quad a.I = a$
0 and 1 is
w.r.t. operation
 $0+a = a + 0$
 $a \cdot 1 = 1 \cdot a$

→ Inverse law corresponds should be
if $a \in B$
then $a+a$

⇒ Properties of

→ Idempotent

$$\begin{aligned} a+a &= a \\ \Rightarrow a+a &= a \\ \Rightarrow (a+a) \cdot 1 &= a \\ \Rightarrow (a+a) \cdot (a+a) &= a \\ \Rightarrow a+a &= a \\ \Rightarrow a &= a \end{aligned}$$

$$\begin{aligned} a \cdot a &= a \\ \Rightarrow (a \cdot a) + &= a \\ \Rightarrow a \cdot a + &= a \\ \Rightarrow a(a+a) &= a \end{aligned}$$

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for multiplication identity element is
 $=1 \quad a \cdot 1 = a \quad 1 = 1$

0 and 1 are identity element
w.r.t. to operator '+' and ' \cdot '.
 $0+a = a + 0 = a \quad \forall a \in B$
 $a \cdot 1 = 1 \cdot a = a \quad \forall a \in B$

→ Inverse law :- every element of set B corresponds both the operation should have inverse element if $a \in B$ then a' is inverse of a then $a+a' = 1$ and $a \cdot a' = 0$

⇒ Properties of Boolean algebras

→ Idempotent law :-

$$\bullet \quad a+a = a$$

$$\Rightarrow a+a = a$$

$$\Rightarrow (a+a) \cdot 1 \quad [\text{Identity law}]$$

$$\Rightarrow (a+a) \cdot (a+a') \quad [\text{Inverse law}]$$

$$\Rightarrow a+a'a' \quad [\text{Distributive law}]$$

$$\Rightarrow a+0 \quad [\text{Identity law}]$$

$$\Rightarrow a$$

• $a \cdot a = a$

$$\Rightarrow (a \cdot a) + 0 \quad [\text{Identity law}]$$

$$\Rightarrow a \cdot a + a a' \quad [\text{Inverse law}]$$

$$\Rightarrow a(a+a') \quad [\text{Distributive law}]$$

$\Rightarrow a \cdot 1$ [Identity law]

$\Rightarrow a$

- $a+1 = 1$
- $a \cdot 0 = 0$

→ Absorption law

- $a + a \cdot b = a$
- $a(a+b) = a$

(Ques 3) In a boolean algebra B prove that Identity elements are complementary to each other if means for $0, 1 \in B$

$$\textcircled{1} \quad 0' = 1 \quad \textcircled{2} \quad 1' = 0$$

$\textcircled{1} \quad 0' + 0 \Rightarrow 1 + 0$ (By Identity law)
 $\Rightarrow 1 //$ Hence proved

$$\textcircled{2} \quad 1' = 0$$

$\Rightarrow 1' \cdot 1$ (By Identity law)
 $\Rightarrow 0 //$ Hence proved

(Ques 2) Simplify following using Boolean algebra

$$\textcircled{i} \quad (a+b)a'b'$$

$$\textcircled{ii} \quad abc + a'b' + c'$$

$$\textcircled{iii} \quad (ab' + c')'$$

$$\textcircled{iv} \quad [a+a'b][a'+ab]$$

$$\textcircled{v} \quad ab + [(a+b').b]'$$

(Ques 8) Let P = it is a simple which describes statements

- (i) $\sim P$ (ii) $P \wedge q$ (iii)
 (v) $P \Rightarrow \sim q$ (vi) $q \vee P$
 (viii) $P \Leftrightarrow \sim q$ (ix)
 (xi) $\sim \sim P$ (xii)

(Ques 4) Let P be "Let q be"
write each form

- (a) Ravi is short
 (b) Ravi is tall
 (c) It is not true
 (d) not handsome

(Ques) Let P be "Let q be"
simple very each of the

$$\textcircled{a} \quad Pvq$$

$$\textcircled{b} \quad P \wedge$$

$$\textcircled{c} \quad \sim(\sim P)$$

Ques 3) Let P = it is cold, q = it is raining
Give a simple verb sentence
which describes each of the following
statements

- (i) $\sim P$ (ii) $P \wedge q$ (iii) $P \vee q$ (iv) $q \leftrightarrow P$
(v) $P \Rightarrow \sim q$ (vi) $q \vee \sim P$ (vii) $\sim P \vee \sim q$
(viii) $P \Leftrightarrow \sim q$ (ix) $\sim \sim q$ (x) $(P \wedge \sim q) \Rightarrow P$
(xi) $\sim \sim P$ (xii) $(P \wedge \sim q) \Rightarrow q$

Ques 4) Let P be "Ravi is tall and
handsome".
Let q be "Ravi is short and
handsome".
Write each statement in the symbolic
form.

- (a) Ravi is short or handsome. $\sim P \vee q$
(b) Ravi is tall or handsome. $P \vee q$
(c) It is not true that Ravi is short or
handsome. $\sim \sim P \vee \sim q$

Ques 5) Let P be "Ravi speaks Tamil" and
 q be "Ravi speaks Hindi". Give a
simple verbal sentence which describes
each of the following

- (a) $P \vee q$ (b) $P \wedge q$ (c) $P \wedge \sim q$ (d) $\sim P \vee \sim q$
(e) $\sim(\sim P)$

De Morgan's law

$$\rightarrow (a+b)' = a'b'$$

$$\rightarrow (ab)' = a'b'$$

This are which quantifier is

Algebra of logic

Quantifier :- when variable is a preposition function. Assume a value, the resulting statement becomes a preposition with truth values

In quantifier we create preposition from a preposition function. It extent to which product is true on a range of element

In english words all, some, many, none and few are used as quantification

Parallel =
Series =

Ques 1) find the function following

(a)

(b)

\Rightarrow There are two types of quantifier soin

$f =$
 $f_1 = a +$

$f_2 = c \cdot$

Ques 2)

Draw
follow

- Existential quantifier or quantification :- which tells us that a predicate is true for every element under the condition
- Universal quantifier or quantification :- which tells us that there is one or more element for predicate is true

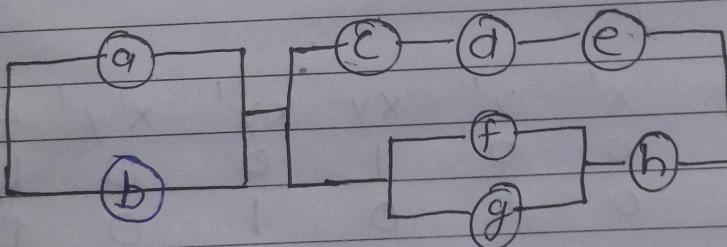
The one which deals with predicates & quantifiers is called as Predicate calculus which

Algebra of electric circuit

$$\text{Parallel} = \text{o/p} \rightarrow x+y$$

$$\text{Series} = \text{o/p} \rightarrow x \cdot y$$

Ques1) find the switching net or switching function or polynomial for the following circuit

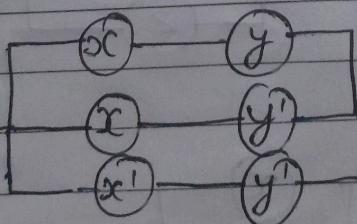


antifler soln $f = f_1 \cdot f_2$

$$f_1 = a+b$$

$$f_2 = cde + (f+g) \cdot h$$

Ques2)



Draw a simple circuit for the following diagram and also verify the

equivalent circuit and the circuit
with the help of truth table

SOLN

$$f = f_1 \cdot f_2 \cdot f_3$$

$$\Rightarrow f_1 = x \cdot y$$

$$\Rightarrow f_2 = xy' \Rightarrow f_3 = x'y'$$

$$\Rightarrow f = (x \cdot y) + (x \cdot y') + (x' \cdot y')$$

$$\Rightarrow x(y+y') + x'y' [DL]$$

$$\Rightarrow x + x'y' [IL]$$

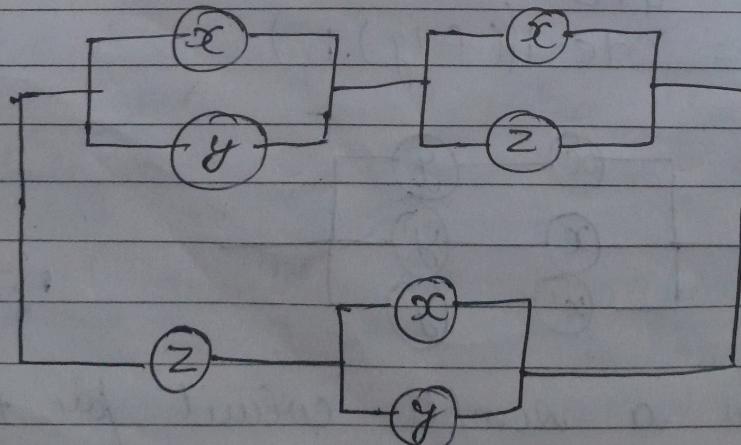
$$\Rightarrow (x+x')(x+y') [DL]$$

$$\Rightarrow 1 \cdot x+y'$$

$$\Rightarrow x+y'$$

x	y	x'	y'	xy	xy'	$x'y'$	f	$x+y$
1	1	0	0	1	0	0	1	1
1	0	0	1	0	1	0	1	1
0	1	1	0	0	0	0	0	0
0	0	1	1	0	0	1	1	1

Aus 3



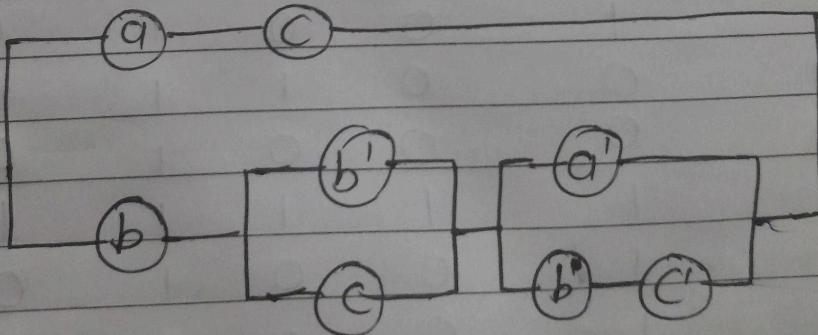
Aus 4)

circuit
table

Date: / / Page no.:

y'	f	x_1y_1'
1	1	0
1	1	0
0	0	1
1	1	1

Ans 4)



$$\underline{\text{Soln}} \quad f_1 = ac$$

$$f_2 = b \cdot (b' + c) \cdot [a' + (bc')]$$

$$f = f_1 + f_2$$

$$= ac + [b \cdot (b' + c) \cdot (a' + (bc'))]$$

$$= ac + [(bb' + bc) \cdot (a' + (bc'))]$$

$$= (ac + bc) \cdot (a' + bc')$$

$$= c(a+b) \cdot (a' + bc')$$

Ans2) PT the
x t +
replaced

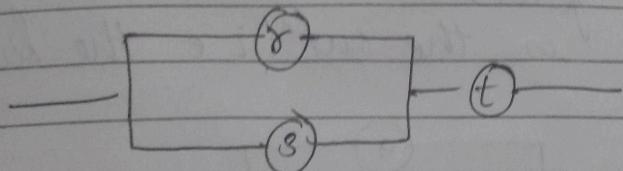
$a \cdot b$	c	a'	b'	c'	f_1	$b' + c$	$b \cdot (b' + c)$	<u>Soln</u>
1	1	1	0	0	0	1	1	x t +
1	1	0	0	0	1	0	0	x t +
1	0	1	0	1	0	1	0	x t +
1	0	0	0	1	1	0	0	x t +
0	1	1	1	0	0	0	1	x t +
0	1	0	1	0	1	0	0	x t +
0	0	1	1	0	0	1	0	x t +
0	0	0	1	1	0	1	0	x t +

Ans2) Draw
expres

bc'	$a' + (bc')$	f_2	f	$a+b$	$c \cdot (a+b)$	Output
0	0	0	1	1	1	0
1	1	0	0	1	0	Design
0	0	0	1	1	1	
0	0	0	0	1	0	
0	1	1	1	1	1	
1	1	0	0	1	0	
0	1	0	0	0	0	
0	1	0	0	0	0	

Ex:- The
doors
close
when
pressed
control

Ques) PT the Boolean function
 $s't + [s(s' + t)\{s' + (st)\}]$ is replaced by the following net



$$\begin{aligned}
 & \text{Soln} \\
 (b') + (c) &= s't + [s(s' + t)\{s' + (st)\}] \\
 &= s't + [ss' + st\{s' + (st)\}] \quad [\text{By DL}] \\
 &= s't + [0 + st\{s' + (st)\}] \quad [aa' = 0] \\
 &= s't + [st + s't] \quad [st \cdot st = st] \\
 &= st + s't \quad [a' + 1 = 1] \\
 &= t(s + s') \\
 & \text{Hence proved}
 \end{aligned}$$

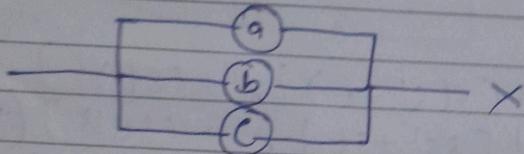
Ques) Draw a simplified circuit of an expression $xy'z + (z + y)x'$

b) Ques) Design of Synthesis Control System

Ex:- The bulb inside a car of two doors is on when one of these door is open and it is also on when a switch of dashboard is pressed draw a diagram for the control path

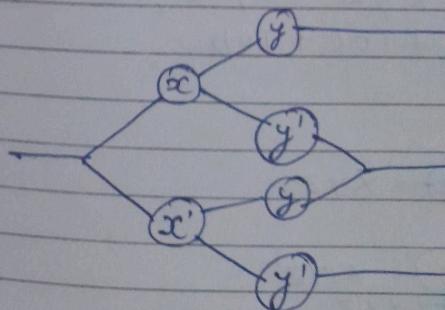
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let $a = 1^{\text{st}}$ door is opened
 $b = 2^{\text{nd}}$ door is opened
 $c = \text{switch of disk board is pressed}$
 $x = \text{is the out i.e. the bulb is on}$



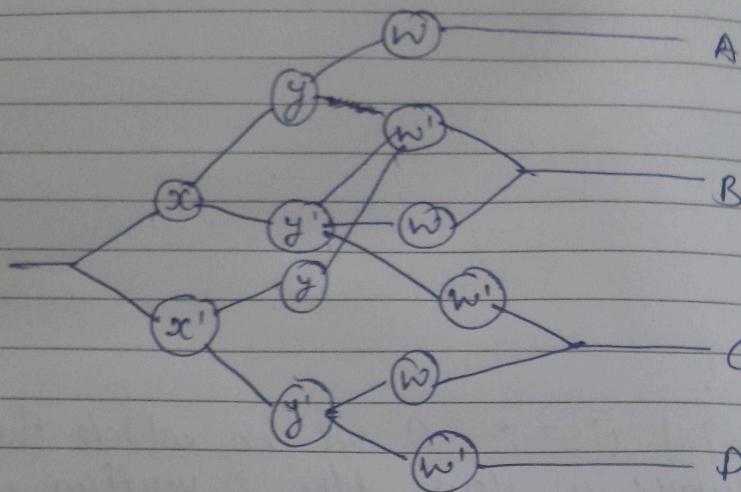
Ex2) A committee has three members # Binomial NET :- A member arrangement is made so that each member has 'Yes' bulb and 'No' button for voting. \Rightarrow Draw a binomial net

A green bulb is on when all are voting and majority is in favour. A red bulb is on when all are voting and majority is in opposition. A yellow bulb is on when all are not voting if any of the member press his vote button and no bulb is on but a bell sounds loudly draw the diagram for its control path



$$\begin{aligned} T_{OA} &= xy & T_{OC} \\ T_{OB} &= xy' + x'y & T = \\ f &= T_{OA} + T_{OB} + T_{OC} \\ &= xy + x'y + xy' + x' \\ &= 1 \end{aligned}$$

→ Draw a binomial net of three variables



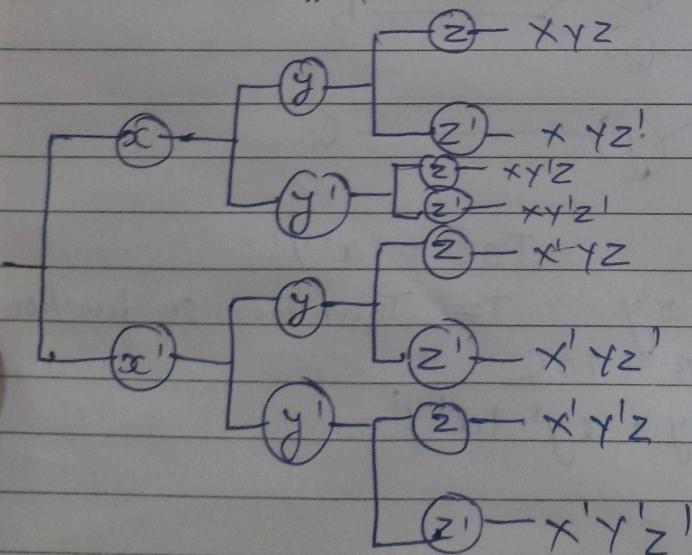
$$T_{OA} = xyA$$

$$T_{OB} = x'y'w + xyw' + x'yw'$$

$$T_{OC} = x'y'w + x'yw' + xy'w', T_{OD} = \cancel{x'y'w'}$$

$$f = T_{OA} + T_{OB} + T_{OC} + T_{OD}$$

Binomial Tree



\Rightarrow Lower Bound. - let $x=50, y=20$

$$\text{then } z_i \cap 50 = \{50, 25, 10, 5, 2, 1\}$$

$$z_j \cap 20 = \{20, 10, 5, 4, 2, 1\}$$

$$z_i \cap z_j = \{10, 5, 2, 1\}$$

\Rightarrow Greatest lower bound = 10 //

\Rightarrow Upper bound = let $x=10, y=4$

$$\text{then } z_i \cap 10 = \{10, 20, 50, 100\}$$

$$z_j \cap 4 = \{4, 20, 100\}$$

$$z_i \cap z_j = \{20, 100\}$$

\Rightarrow least upper bound = 20 //

\Rightarrow ^{Upper} ~~Lower~~ Bound = let $x=25, y=10$

$$\text{then } z_i \cap 25 = \{25, 50, 100\}$$

$$z_j \cap 10 = \{10, 50, 20, 100\}$$

$$z_i \cap z_j = \{50, 100\}$$

\Rightarrow ~~Greatest~~ least ^{upper} ~~lower~~ bound = 50

\Rightarrow ^{lower} ~~upper~~ bound = let $x=25, y=10$

$$z_i \cap 25 = \{25, 5, 1\}$$

$$z_j \cap 10 = \{10, 5, 2, 1\}$$

$$z_i \cap z_j = \{5, 1\}$$

Greater lower bound 5

- ### # Properties
- ① A set m bound or
 - ② If the are uniq
 - ③ If for elem is call

- Q) Draw a
- x =
 - x ≤ y
 - (i) 10B
 - (ii) the 10
 - (iii) 10B

Properties

Date: / / Page no. / /

- ① A set may have no upper & lower bound or it may have many
- ② If the meet and join exist they are unique
- ③ If a Poset has GLB and LUB for every pair elements then it is called lattice

Q) Draw a Hasse diagram for (X, \leq)

$$X = \{2, 3, 6, 12, 24, 36\} \text{ and}$$

$X \leq Y$ if X/Y find

i) LUB and GLB of $A = \{2, 3, 4\}$

ii) the LUB and GLB of $B = \{2, 3\}$

iii) LUB and the GLB of $C = \{6, 12\}$

20

5, 2, 13
4, 2, 13

$y = 4$

50, 100
0 { }

$y = 10$

{ }

0 { }

50

10

Join Semi
In a poset
exist for
then pos
lattice
 \Rightarrow

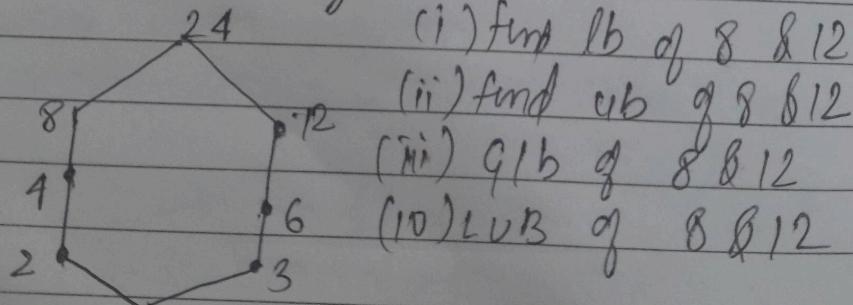
is not the
not their

Total are
partial
s are

(i) compare
(ii) incompara

If
compar
Total
a few

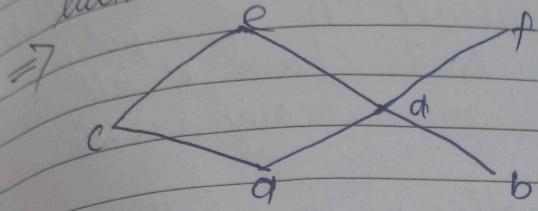
Q2) for hasse diagram



- (i) find Lb of 8 & 12
- (ii) find Ub of 8 & 12
- (iii) GLb of 8 & 12
- (iv) LUB of 8 & 12

\Rightarrow Well a
a we
poset
orderin

Join Semi Lattice
 In a poset if lub/join/supremum/ \vee_0 ,
 exist for every pair of element
 then poset is called join semi
 lattice



it is neither
 join and meet
 semi lattice

as join of $e \wedge f$
 is not their and meet of $a \wedge b$ is
 not their

01/03/2023

Total ordering :- S is a set with
 partial ordering so the elements of
 S are

- (i) Comparable $\rightarrow (3, 9)$
- (ii) incomparable $\rightarrow (5, 7)$

If all the pairs in relation are
 comparable then it is known as
 Total ordering provided set S is
 a poset

\Rightarrow Well ordered sets (S, \leq) this is
 a well ordered set if it is a
 poset such that ' \leq ' is a total
 ordering and every non empty subset

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of S has atleast one element

Chains & Anti chains :- A subset ' A ' of σ poset is called a Chain if all the elements of ' A ' are comparable. A subset ' A ' of a poset is called an Anti chain if all the elements of ' A ' are incomparable.

Graph

↪ It is GC V, E

↪ vertex/nod

↪ Edge/line

Distributed Lattice:- The lattice ' P ' denoted by (P, \wedge, \vee) is distributed lattice if it hold distributed law

↪ degree of

↪ this is

↪ every graph

↪ degree

↪ degree of

$$a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$$

and $a \vee (b \wedge c) = (a \vee b) \wedge (a \vee c)$

$\forall a, b, c \in A$

Compliment lattice :- Let (L, \wedge, \vee) is a lattice and $0, 1 \in L$ such that $0 \leq a \leq 1 \quad \forall a \in L$ then

⇒ Graph :-

cl b

$v = \{$

v 's el

point)

called

$$a \vee 1 = a, a \wedge 1 = a$$
$$a \wedge 0 = 0, a \vee 0 = a$$

Now if for all $a \in L$ there is $a' \in L$ such that $a \wedge a' = 0, a \vee a' = 1$ then a' is called 'complement' of a and lattice is called complement lattice

→ Self

a

when same

lement

subset 'A'

chain if

we compareable
of a poset
if all the
variable

lattice 'P'

distributed
ubuted

v(anc)

n(avc)

E A

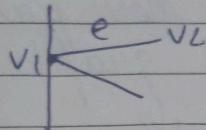
(n, v) is
ch thata' E C
a' = 1element
called

Graph

It is collection of vertex & edges
 $G(V, E)$

vertex/nodes/point = $\{v_1, v_2, \dots, v_n\}$

Edge/Line/Branch = $\{e_1, e_2, \dots, e_n\}$



degree of vertex = total no. of edges

this is undirected graph

Every tree is a graph but

every graph may or may not be a tree

degree of root = 0 (in tree)

degree of vertex = 1 (in tree)

\Rightarrow Graph :- A graph $G(V, E)$ consist of

a set of objects

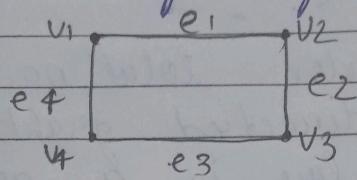
$$V = \{v_1, \dots, v_n\}, E = \{e_1, \dots, e_n\}$$

v's element are called vertices (or
point) Node) and E's elements are
called edges (lines/branches)

\Rightarrow Self Loop :- An edge is said to be
a self loop (or simply a loop)
when its both the vertices are
same.

→ Parallel edges :- If there are two or more than two edges having the same pair of vertices then such edges are called parallel edges or multiple edges

→ Simple graph :- A graph that has neither self loop nor parallel edges is called simple graph



→ Multi / Pseudo graph :- A graph that has self loop or parallel edges is called multi / pseudo graph

→ Incident edge :- Let v_i be an end vertex of some edge e_j then we can say that edge e_j is incident on the vertex v_i .
for example :- e_1 , e_3 and e_4 are incident on v_2

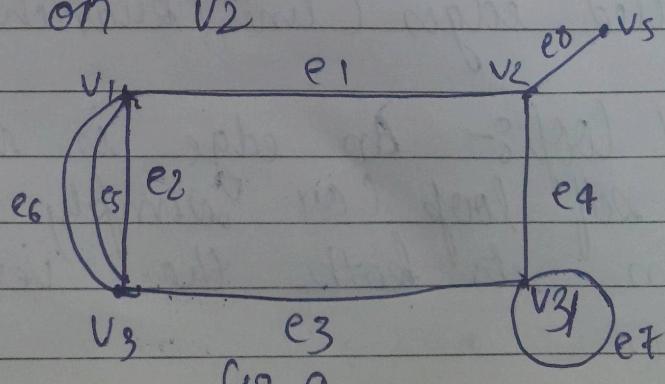


fig 9

→ Adjacent edges are called adjacent edges

→ Adjacent vertices are called adjacent vertices

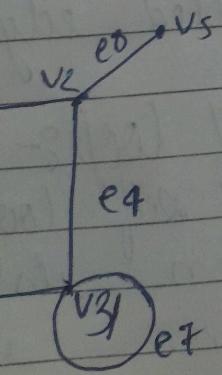
→ Degree of a vertex
The degree of a vertex is the number of edges incident on it.

Page no. _____ Date: / / Page no. _____
There are two or more edges having the same vertices then called parallel edges

that has parallel edges graph

A graph that has parallel edges graph

Let v_i be an end vertex of edge e_j then edge e_j is incident on v_i and e_j are



→ Adjacent edges :- Two non parallel edges are called adjacent if they are incident on common vertex

→ Adjacent vertices :- Two vertices are called adjacent if there is an edge joining them. It means v_2 and v_4 are adjacent vertices (from fig a)

→ Degree of vertices :-

The degree of vertices v_i is denoted by $\text{degree } v_i$ (see fig a)

$$\begin{aligned}\text{degree } (v_1) &= 4 \quad \text{find} \\ \text{degree } (v_4) &= 4 \quad \text{same for all}\end{aligned}$$

Adjacent edges :- Two non parallel edges and call adjacent if they are incident on common vertex.

Adjacent vertices :- Two vertices are called adjacent if there is an edge joining them it means v_2 and v_1 are adjacent vertices (from fig a)

Degree of vertices :-

The degree of vertices v_i is denoted by degree v_i (see fig a)

$$\text{degree } (v_1) = 4$$

$$\text{degree } (v_4) = 4 \quad \begin{matrix} \text{find} \\ \text{same for all} \end{matrix}$$

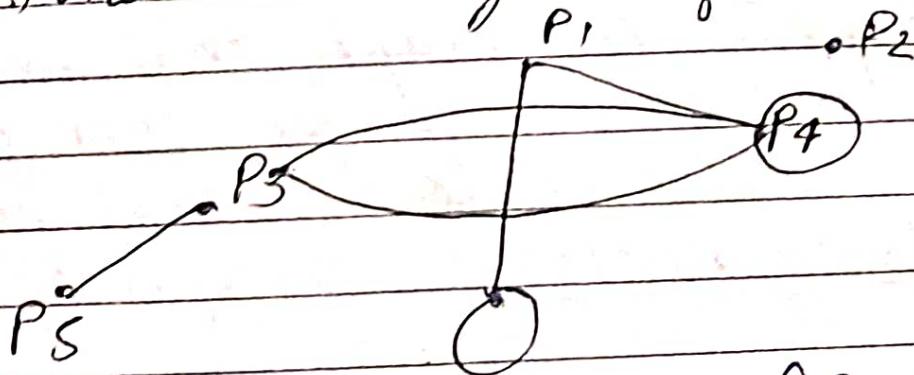
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Regular graph :- A graph G in which all vertices are of equal degree is called regular graph.

ex:-



Find the degree of each vertex



$$\deg (P_1) = 2 \quad \deg (P_2) = 0, \deg (P_3) = 3$$

$$\deg (P_4) = 3 \quad \deg (P_5) = 1, \deg (P_6) = 3$$

→ Null graph :- The null graph is a graph containing no edges, the null graph with n vertices is denoted by N_n .
ex:- $\bullet N_1$ $\bullet N_2$ $\bullet N_3$

degree is 0 in null graph
for each vertex

→ Isolated vertex :- A vertex with 0 degree is called isolated vertex.

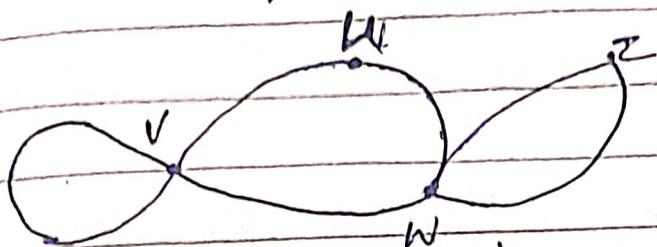
→ Pendant or Peivial graph :- A vertex of degree one is called pendant graph.

→ finite & infinite graph :- A graph $G(V, E)$ is finite if both V & E are finite. The graph $G(V, E)$ is infinite if both V & E are infinite.

→ Connected graph :- A graph that is in one piece is said to be connected whereas one split into several pieces is called disconnected graph.

→ Sub graph :- Subgraph of a graph G is a graph whose all the

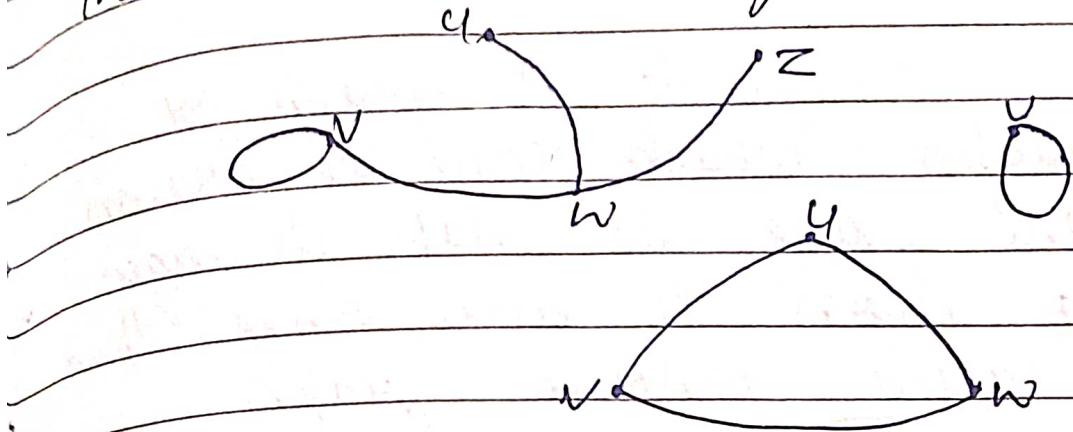
vertices belongs to $V(G)$ and all whose
all the edges belongs to $E(G)$
for example if G is the connected
graph



$$\text{where } V(G) = \{v, u, w, z\}$$

$$E(G) = \{vu, vw, vw, wz, wz\}$$

then the following graphs are subgraphs



→ Cycle graph :- A cycle graph is a graph of single cycle. The cycle graph with n vertices is denoted

by C_n

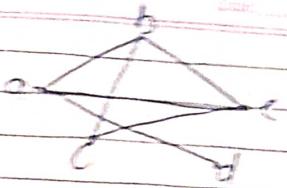
ex:- C_1



C_3

Handshaking Theorem

Consider any graph. The sum of all the degrees of a graph is equal to twice the number of edges.



Face vertex $\Rightarrow \{a, b, c, d, e\}$

edges = $\{ab, ac, ad, bc, be, cd, ce\}$

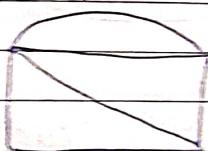
$\deg(a) = 3, \deg(b) = 3, \deg(c) = 2$
 $\deg(d) = 2, \deg(e) = 4$

is defined
path ! that
the short

eccentricity
vertex
distance
first
metrically
when

→ Planar Graph & A graph or "multigraph" which can be drawn in the plane so that its edges do not cross to each other. Then it is called PLANAR Graph.

ex :-

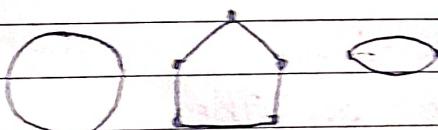


→ Centri-
minimum
as the

$E(a)$ =
 $E(d)$ =
 there

→ Circuit \Rightarrow The closed walk in which no vertex appears once is called circuit

ex:-



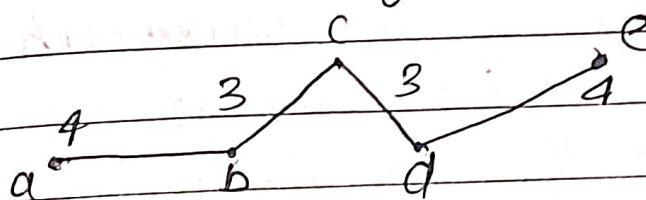
→ In G
 G the
 of G
 G_{cu}

→ Distance - The distance $d(v_i, v_j)$ b/w vertex any given pair of vertex $v_i \& v_j$ then

is defined to the length of shortest path that is the no. of edges in the shortest path

→ Eccentricity:- The eccentricity e of a vertex v is defined as its maximum distance from any other vertex furthest from v this can be mathematically as $e(v) = \max d(v_i, v_j)$ where $i = 1, 2, \dots, n$.

→ Centre:- A vertex in a graph G with minimum eccentricity is referred to as the center of G



$$e(a) = 4, e(b) = 3, e(c) = 2$$

$$e(d) = 3, e(e) = 4.$$

these are the eccentricity values

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→ In & Out degree :- In directed graph G the out degree of vertex v of G is denoted by $\deg_G^+(v)$ or $\deg_G(v)$, if v is the no. of edges in the directed graph G of the vertex v beginning at v and then in degrees of vertex is denoted

If the no. of edges ending at v the sum of in and out degrees of a vertex is called the total degree of the vertex.

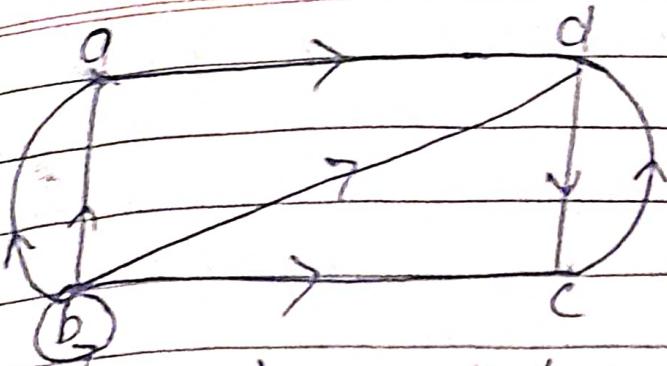
The vertex with '0' in degree is called source and a vertex with '0' out degree is called sink.

If $G(V, E)$ be a directed graph with e edges

$$\sum_{v \in V} \deg_q^+(v) = \sum_{v \in V} \deg_q(v)$$

It means the sum of out degrees of the vertices of a digraph equals the sum of in degrees of the vertices which equals the number of edges in G .

Proof:- Any directed edge (u, v) contributes 1 to the degree of v and 1 to the out degree of u . The loop at v contributes 1 to the in degree and 1 to the out degree of v hence proved.



$\Rightarrow \text{Indegree}(a) = 2 \Rightarrow \text{Indegree}(b) = 1$

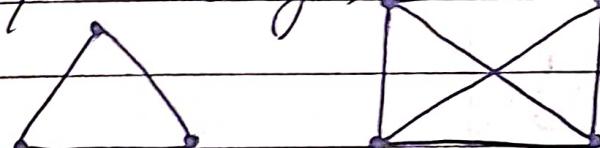
$\text{outdegree}(a) = 1, \text{outdegree}(b) = 5$

$\Rightarrow \text{Indegree}(c) = 2 \Rightarrow \text{Indegree}(d) = 3$

$\Rightarrow \text{outdegree}(c) = 1, \text{outdegree}(d) = 1$

\Rightarrow Complete Graph :- A simple graph G is said to be complete if every vertex in G is connected with every other vertex. If G contains exactly 1 edge b/w each pair of distinct vertices a complete graph usually denoted by ' K_n ' and it should be noted that K_n has exactly $n(n-1)/2$ edges.

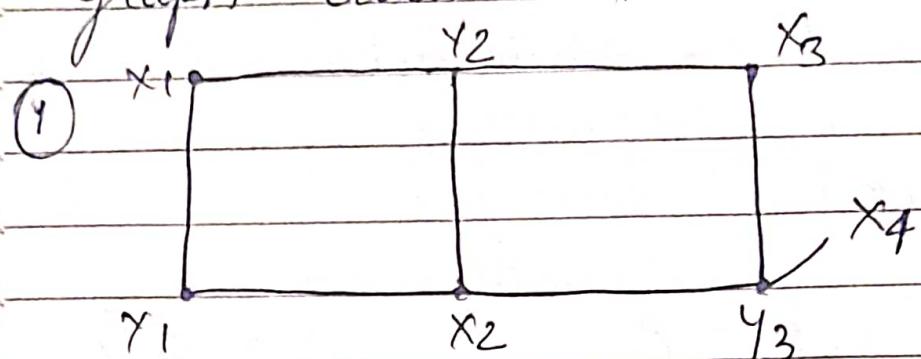
ex:-



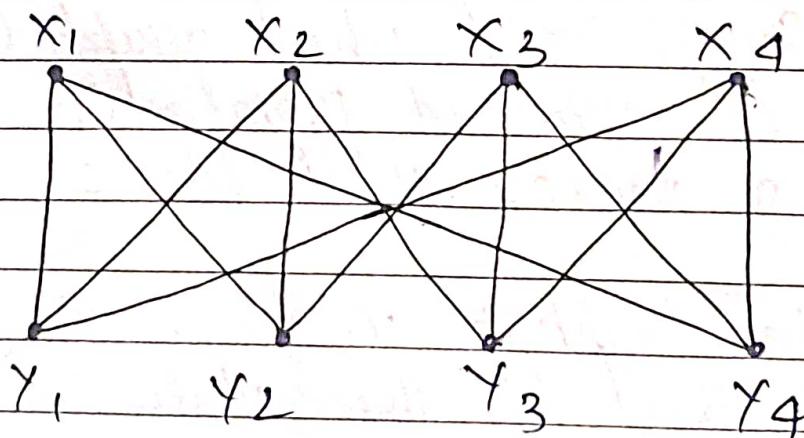
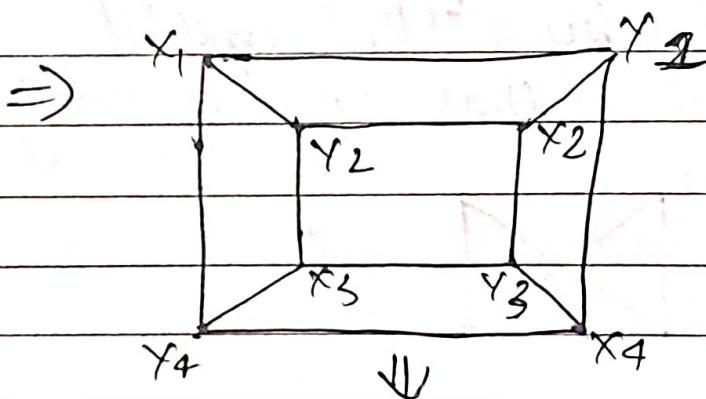
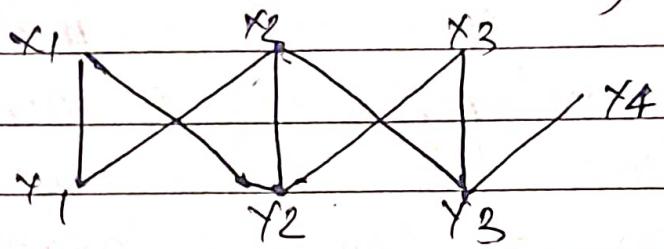
\Rightarrow Wheel Graph :- A wheel graph (W_n) ($n > 3$) is obtained from C_n by adding a vertex v inside C .

\Rightarrow Bipartite Graph :- A graph $G(V, E)$ is Bipartite if the vertex set V

can be partitioned on the two subsets (V_1 & V_2) such that every edge in E connects a vertex in V_1 and a vertex V_2 , V_1 and V_2 is called a bipartition in G . A Bipartite graph doesn't have a loop.



↓ converting it into Bipartite



→ Complete Bipartite: The complete Bipartite graph on m & n vertices denoted by $K_{m,n}$. is a graph whose vertex set is partition into sets V_1 with m vertices & V_2 with n vertices in which there is an edge b/w each pair of vertices $v_1 \& v_2$ where v_1 is in V_1 and v_2 is in V_2

⇒ Any graph in the form of $K_{1,n}$ is called star graph.

⇒ A complete Bipartite graph $K_{m,n}$ is not regular graph